

# Secure Protocols via AI and CPSA

Lauren Brandt

Mentors: Dr. Joshua Guttman and Dr. Andres Molina-Markham

Research Program Lead: Gabby Raymond

HCSS

May 13, 2025

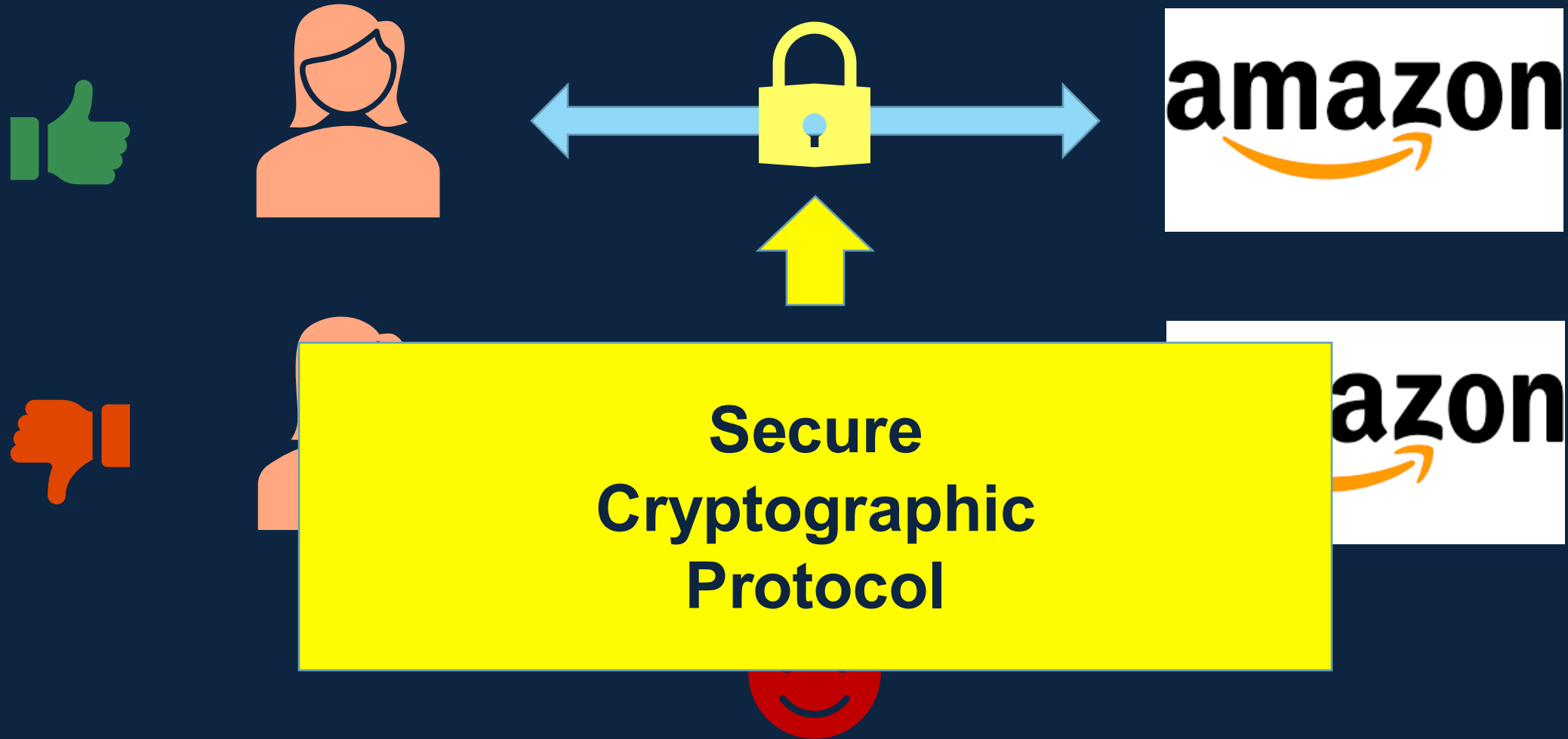
# Overview

## Improving the usability of CPSA. Combining formal methods with LLMs.

- **LLMs for CPSA:** LLMs act as translators and assistants, making CPSA more user-friendly for both experts and novices.
- **CPSA for LLMs:** Risk minimal: CPSA verifies the LLM outputs
- **Method Highlights:** Natural language (NL) format for specifying protocols, reliable and automatic scoring method

LLM – Large Language Model  
CPSA – Cryptographic Protocol Shapes Analyzer

# Cryptographic Protocol



# Hard to Design Protocols

## SSH

24 Draft Versions  
9 Years

## IPsec

10 Draft Versions  
3 Years

## TLS 1.3

28 Draft Versions  
4 Years

## DTLS 1.3

43 Draft Versions  
6 Years

**Long and error prone process to develop  
these secure protocol standards**

# Are they using any tools to create these protocols? Yes.



Tamarin

Can LLMs help make these tools more usable? Maybe.



Isabelle



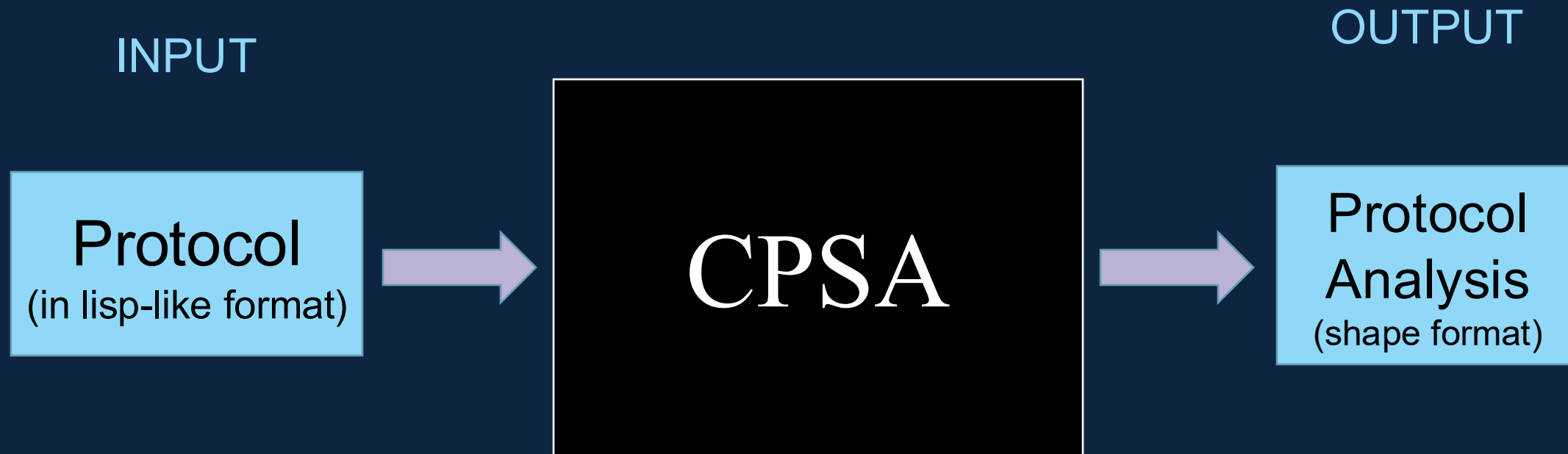
ProVerif

**CPSA**

## But they can be very hard to use.

(They can be so difficult to use, in fact, the international formal methods community has created a working group that aims to make these tools more usable)

# Cryptographic Protocol Shapes Analyzer (CPSA)



Shapes in the CPSA analysis reveals authentication and confidentiality properties about the protocol.

Shapes = all essentially different runs of a protocol in a strong adversary environment

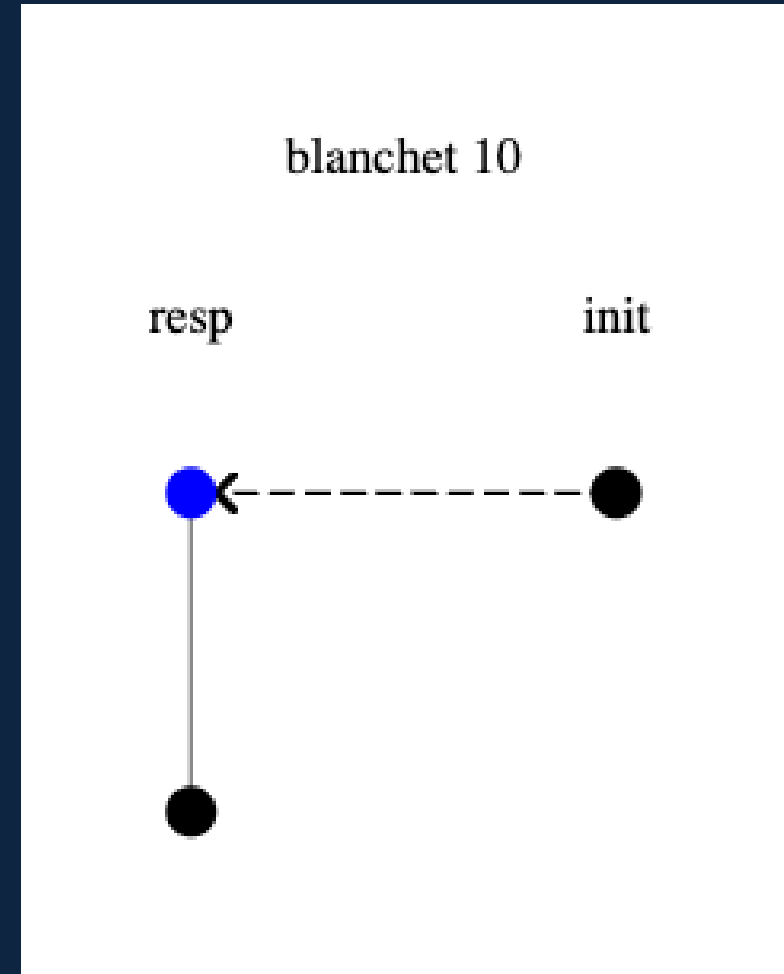
# Cryptographic Protocol Shapes Analyzer (CPSA)

## INPUT

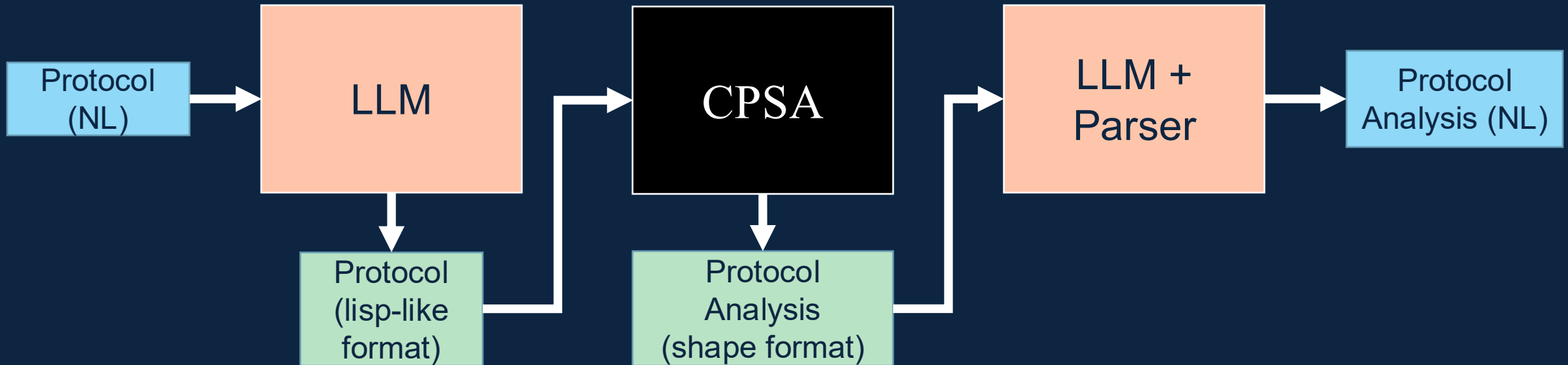
```
(defprotocol blanchet basic
  (defrole init
    (vars (a b name) (s skey) (d text))
    (trace
      (send (enc (enc s (privk a)) (pubk b)))
      (recv (enc d s))))
  (defrole resp
    (vars (a b name) (s skey) (d text))
    (trace
      (recv (enc (enc s (privk a)) (pubk b)))
      (send (enc d s)))))
```

These can be hard  
to create and understand  
Can LLMs Help?

## OUTPUT



# Technical Approach



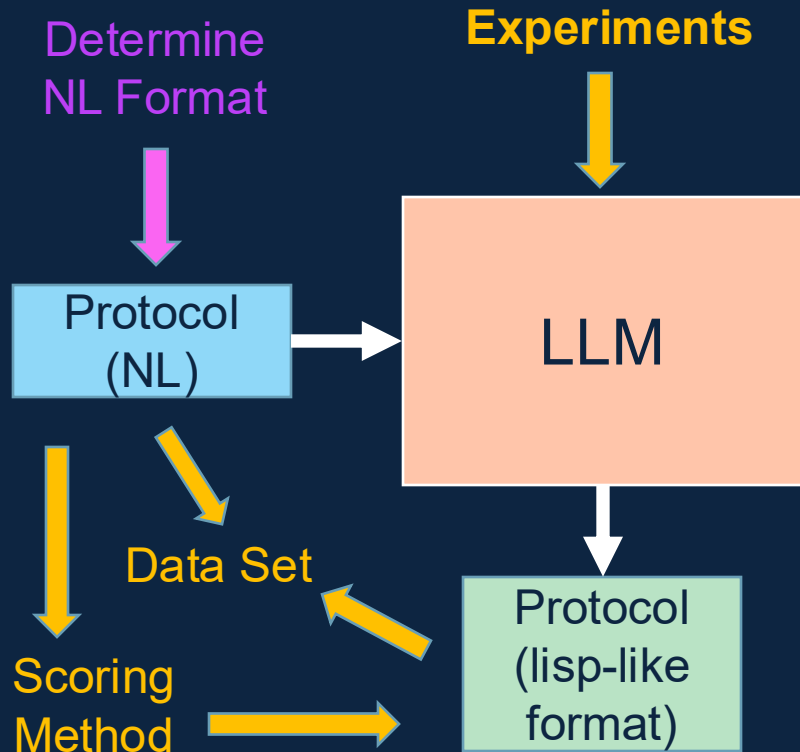
NL – Natural Language

LLM – Large Language Model

CPSA – Cryptographic Protocol Shapes Analyzer

LLMs acting as a translator between natural language and CPSA input/output formats.

# NL → CPSA LLM



## Main Questions:

1. How effective are LLMs at this translation?
2. How does this benefit a CPSA user?
3. How do we handle hallucinations?

# The NL Format

Basic overview  
of protocol



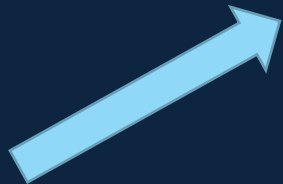
"The Blanchet protocol is a two-party protocol that allows two parties, Alice and Bob, to exchange a shared key, '(s skey)', and a secret message, '(d data)', using this shared key.

Number of  
messages



There are only 2 messages total, so only 2 sends and receives should be in the CPSA protocol. Alice should have one send and one recv in her defrole and Bob should have one recv and one send in his defrole.

Clear description  
of protocol from  
**each role's POV**,  
(with careful wording for  
the double encryption.)



Here is how it works from Alice's point of view: Alice sends the first message and it contains an encryption within an encryption. In this message, Alice encrypts both a shared key 's' and Bob's name 'b' with her private key '(privk a)', and then this encrypted message is further encrypted with Bob's public key '(pubk b)'. Next, Alice recvs a piece of data 'd' encrypted with the shared-key 's'.



Here is how it works from Bob's point of view:  
Bob receives a message that contains an encryption within an encryption. In this message, a shared key 's' and Bob's name 'b' is encrypted with Alice's private key '(privk a)', and then this encrypted message is further encrypted with Bob's public key '(pubk b)'. Next, Bob encrypt a piece of data, 'd', with the shared key that Alice sent in message one, 's'. Bob sends this encrypted piece of data."

# The NL Format – Zooming In

Here is how it works from Alices' point of view:

Alice sends the first message and it contains an encryption within an encryption.

In this message, Alice encrypts both a shared key 's' and Bob's name 'b' with her private key '(privk a)'. This encrypted message is further encrypted with Bob's public key '(pubk b)'.

Next, Alice receives a piece of data 'd' encrypted with the shared-key 's'.

**If the user understands the protocol,  
the NL format should not be difficult to write.**

## Scoring Method

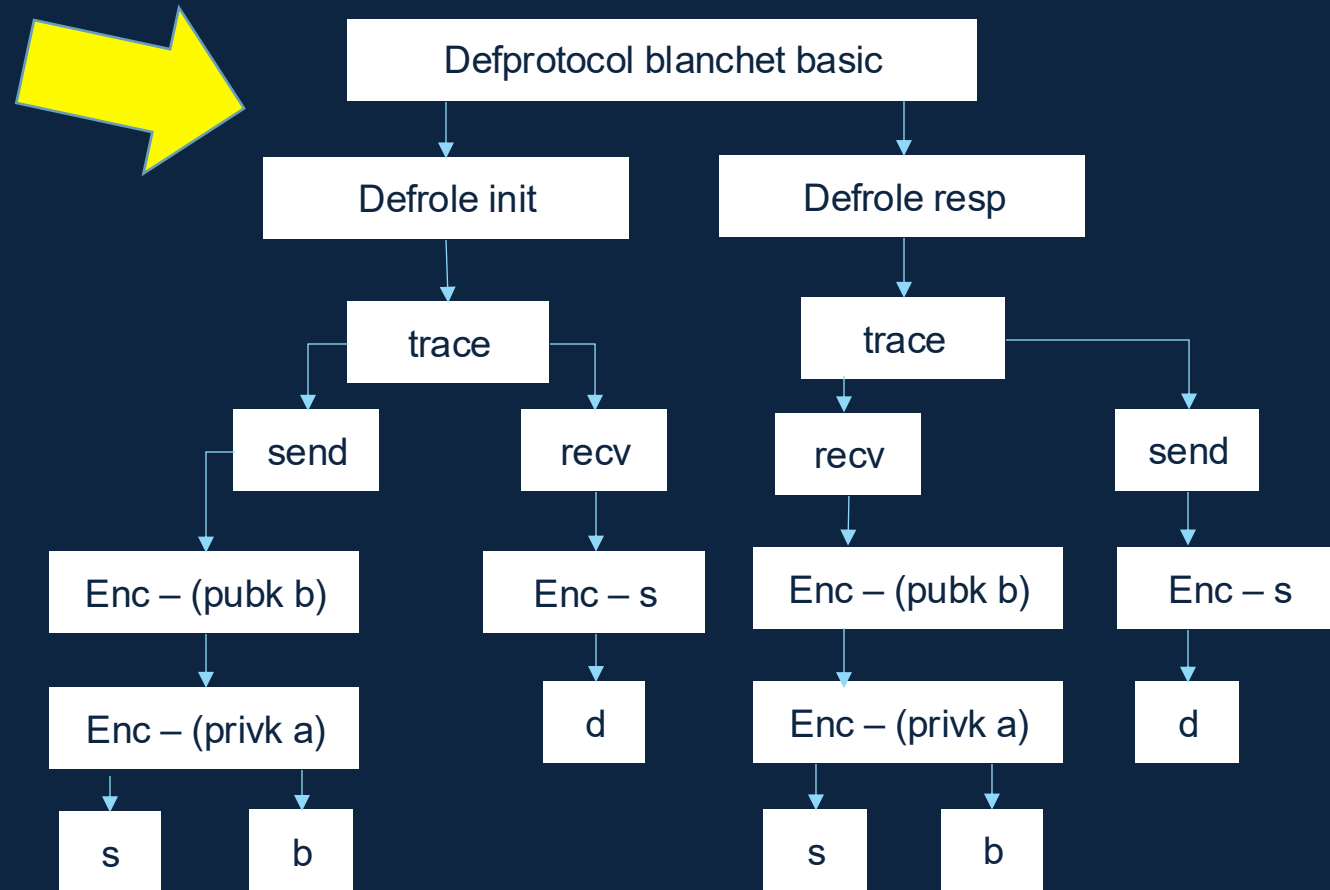
- **Goal:** Evaluate how accurately a LLM translates protocol messages.
- First concerned about **semantics**:
  - LLMs already match the CPSA formatting very closely; CPSA to help with any syntax issues
  - This process is completely unhelpful if the generated protocols run through CPSA without error, but the protocol is completely different from the user's intent.
  - Long-term concerned about both semantics and syntax
- **Idea:** Decompose CPSA into trees and calculate the tree edit distance between actual and generated CPSA

```

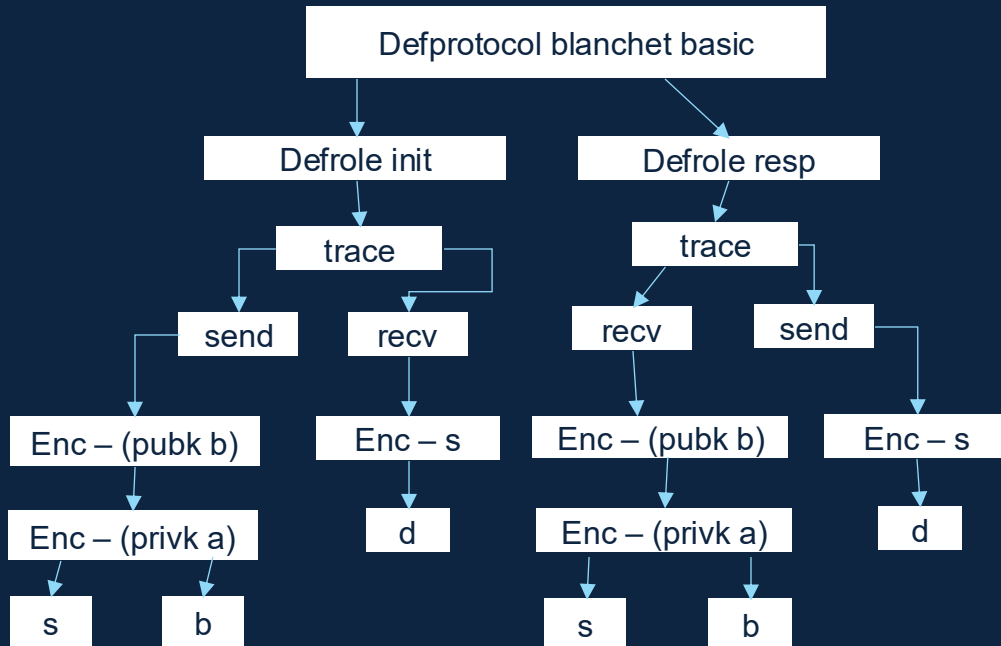
(defprotocol blanchet-fixed basic
  (defrole init
    (vars (a b name) (s skey) (d data))
    (trace
      (send (enc (enc s b (privk a)) (pubk b)))
      (recv (enc d s))))
  (defrole resp
    (vars (a b name) (s skey) (d data))
    (trace
      (recv (enc (enc s b (privk a)) (pubk b)))
      (send (enc d s)))))

```

**CPSA can be  
decomposed into trees.**

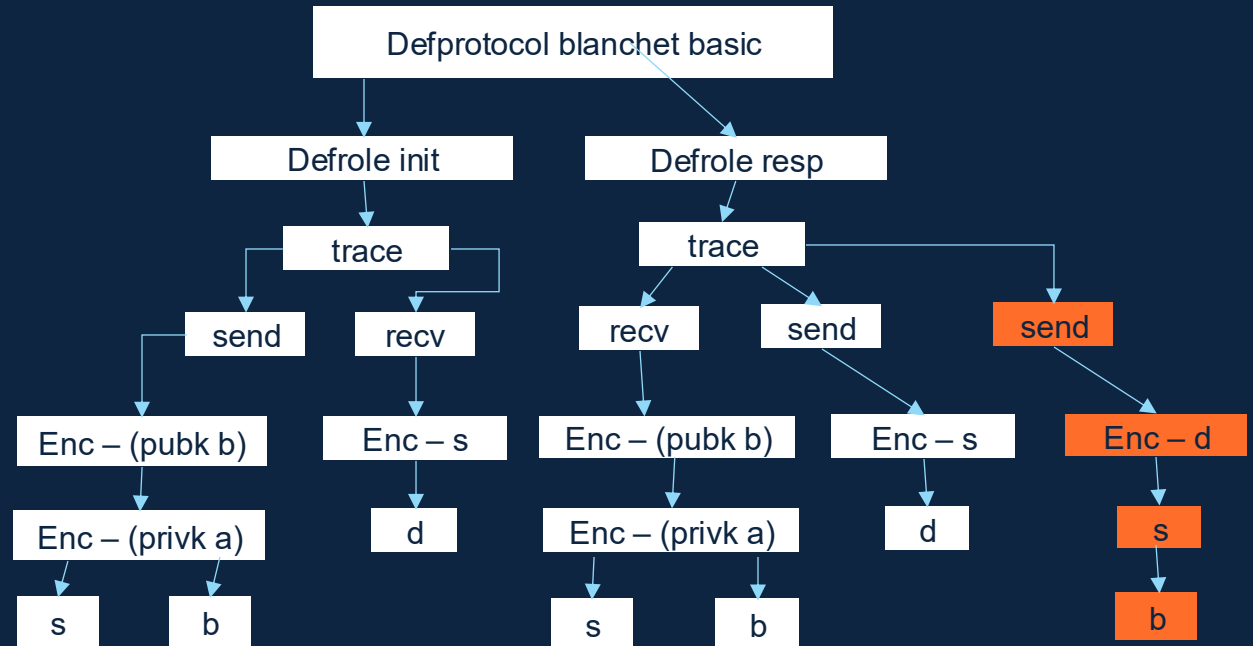


## Actual CPSA Tree



Full tree cost: 21

## Generated CPSA Tree



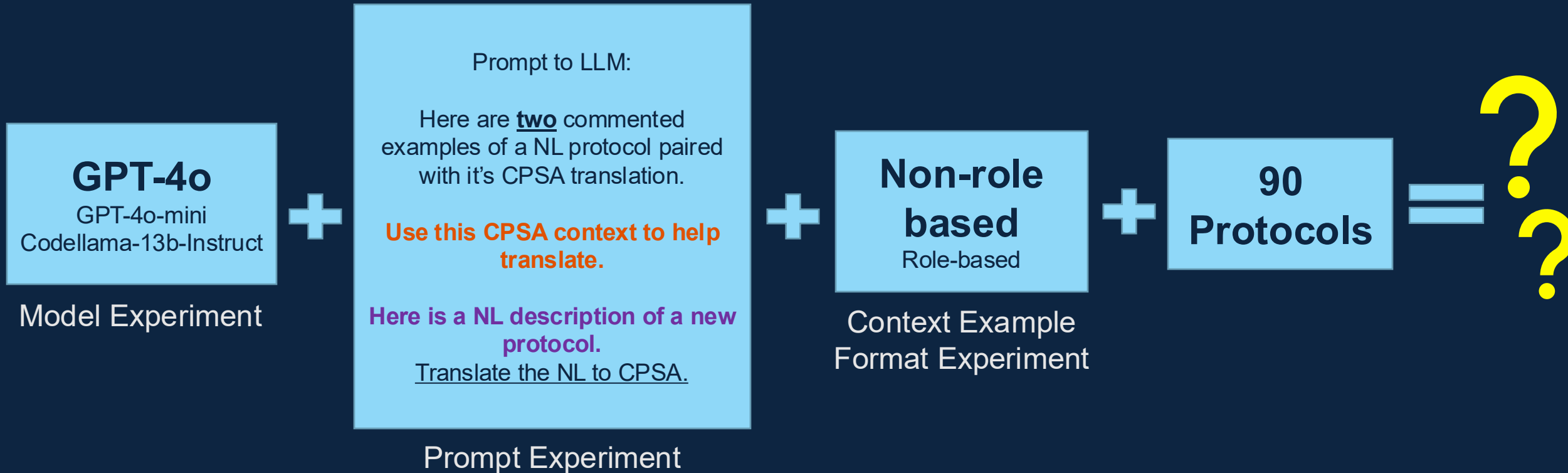
Tree edit distance: 4

Find the tree edit distance between actual and generated CPSA trees.

Tree edit distance: Minimum number of additions, deletions, or modifications to turn one tree into another

# Experiments

## LAST

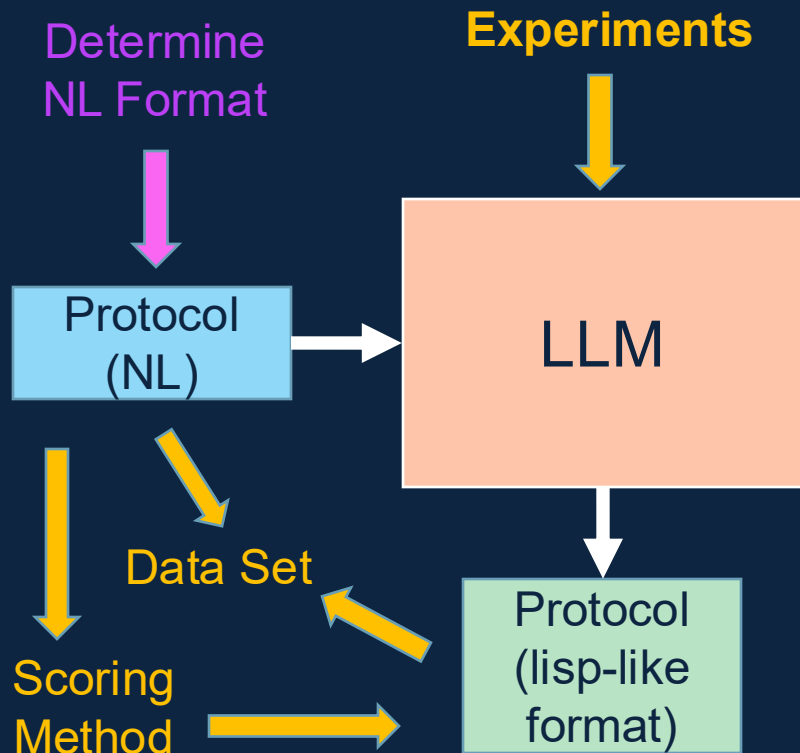


### Translation Accuracy

### Translation Accuracy (%)

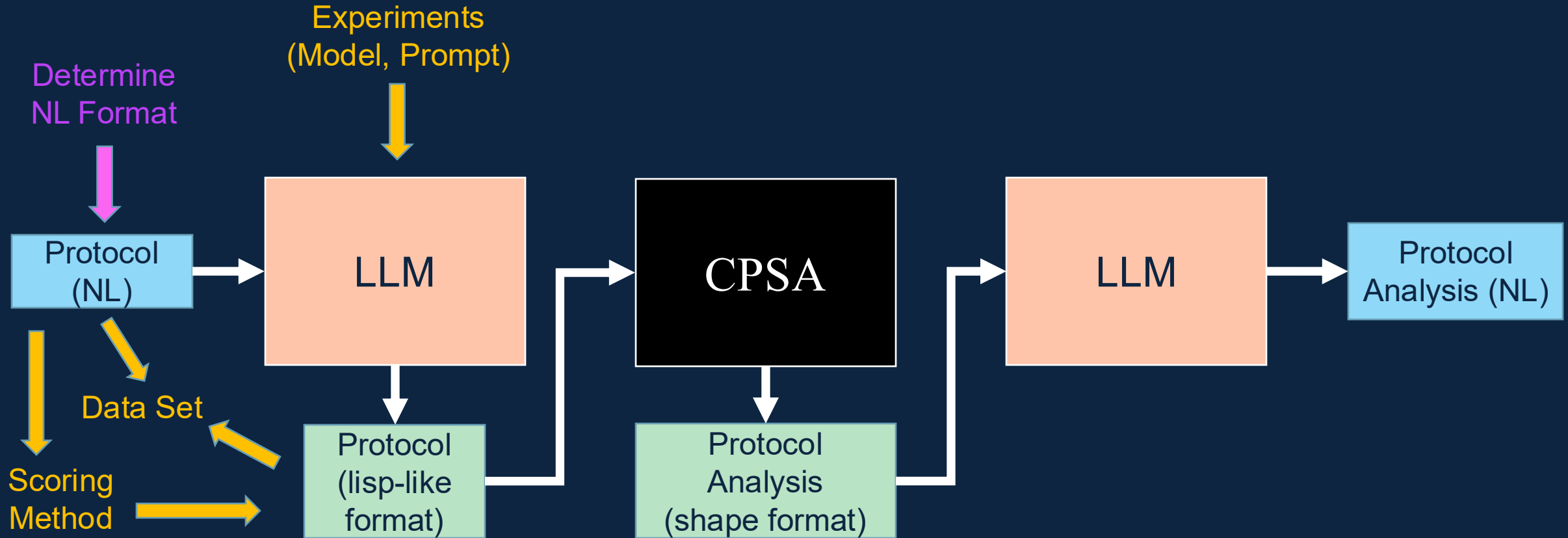


# NL → CPSA Summary

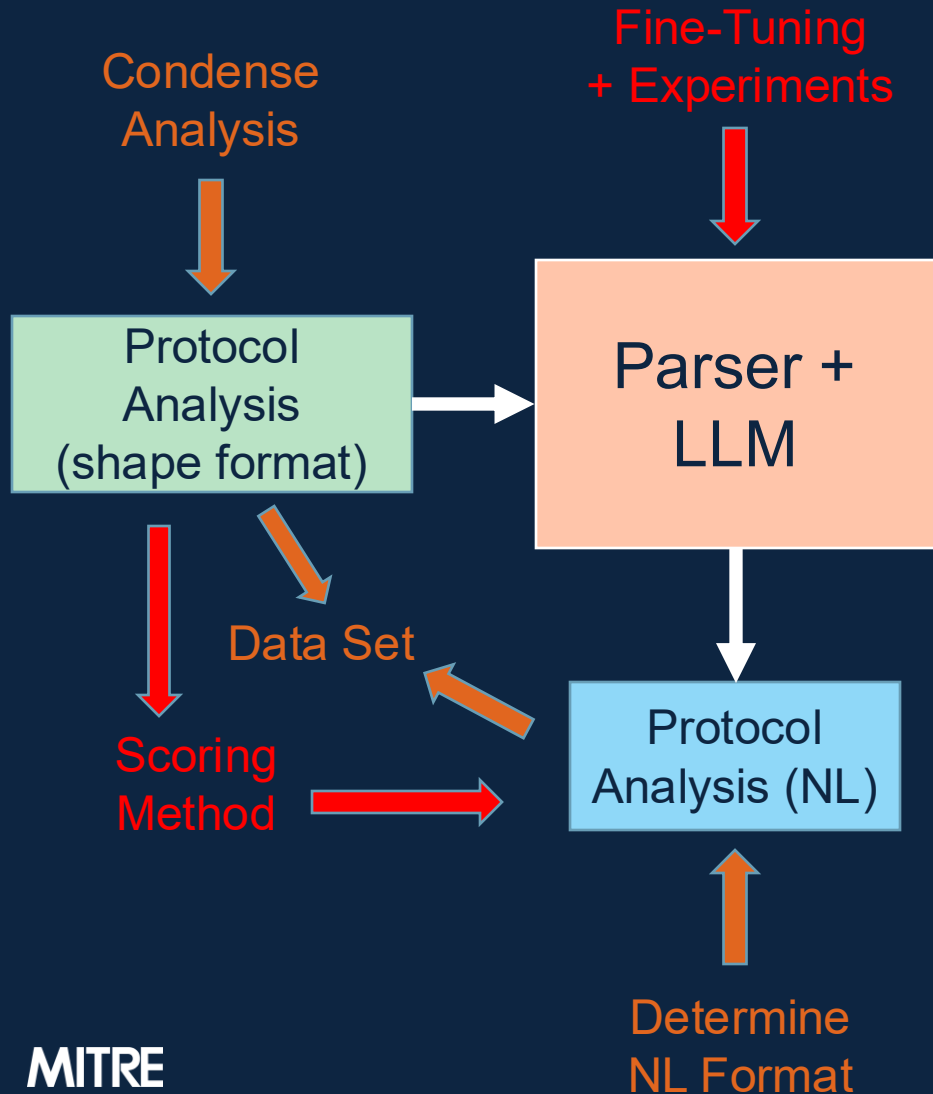


1. How effective are LLMs at this translation?
  - Very effective – 98.3% translation accuracy
2. How does this benefit a CPSA user?
  - For the novice: Easier to learn and use CPSA
  - For the expert: Significantly minimizes manual effort
3. How do we handle hallucinations?
  - Accept the error – new protocol
  - Trace down and fix the error – which we already do

# Research Plan Progress



# CPSA → NL Main Question



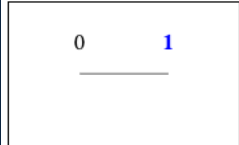
How can we use LLMs and other techniques to:

- identify protocol weaknesses
- strengthen them

# Determine NL Format

```
(comment "CPSA 4.4.2")
(comment "Extracted shapes")
(comment "CPSA 4.4.2")
(comment "All input read from paywave_1.scm")
```

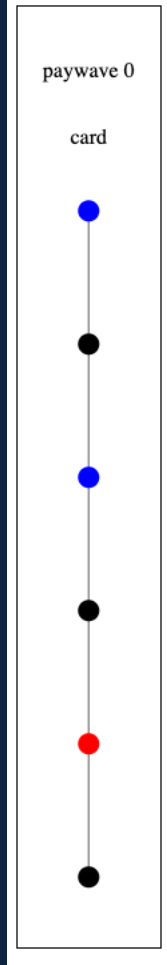
Tree 0, POV 0.



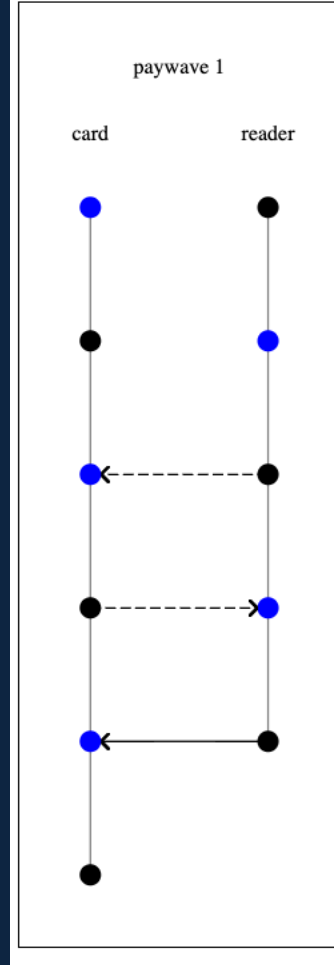
```
(defprotocol paywave basic
  (defrole reader
    (vars (r c name) (opts nr text) (sel nc mesg))
    (trace (send (cat r opts)) (recv (cat r c sel)) (send (cat r c nr))
      (recv (cat r c nr nc)) (send (enc nr nc (privk r)))
      (recv (enc nr nc (privk c))))
    (uniq-orig opts nr))
  (defrole card
    (vars (r c name) (opts nr mesg) (sel nc text))
    (trace (recv (cat r opts)) (send (cat r c sel)) (recv (cat r c nr))
      (send (cat r c nr nc)) (recv (enc nr nc (privk r)))
      (send (enc nr nc (privk c))))
    (uniq-orig sel nc))
  (defgenrule neqRl_indx
    (forall ((x indx)) (implies (fact neq x x) (false))))
  (defgenrule neqRl_strd
    (forall ((x strd)) (implies (fact neq x x) (false))))
  (defgenrule neqRl_mesg
    (forall ((x mesg)) (implies (fact neq x x) (false)))))
```

```
(defskelton paywave
  (vars (opts nr mesg) (sel nc text) (r c name))
  (defstrand card 6 (opts opts) (nr nr) (sel sel) (nc nc) (r r) (c c))
  (non-orig (privk r) (privk c))
  (uniq-orig sel nc)
  (traces
    ((recv (cat r opts)) (send (cat r c sel)) (recv (cat r c nr))
      (send (cat r c nr nc)) (recv (enc nr nc (privk r)))
      (send (enc nr nc (privk c)))))
  (label 0)
  (unrealized (0 4))
  (origs (sel (0 1)) (nc (0 3)))
  (comment "1 in cohort - 1 not yet seen"))
```

Item 0, Child: 1.



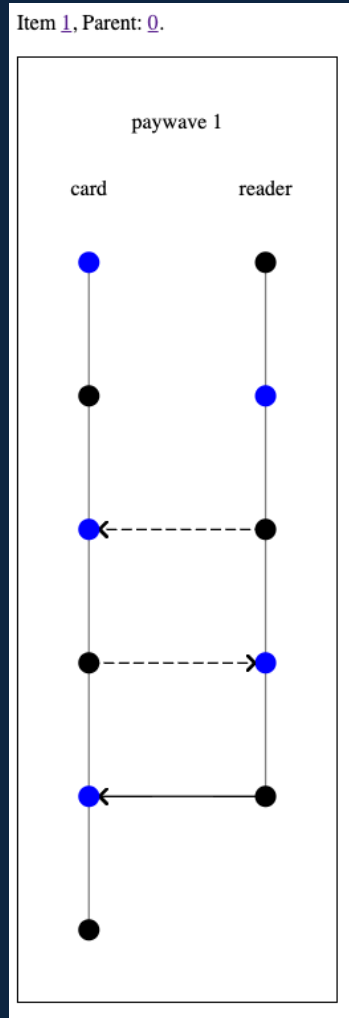
Item 1, Parent: 0.



Some of the information that CPSA gives you for just one query.

```
(defskelton paywave
  (vars (opts sel mesg) (sel-0 nc opts-0 nr text) (r c c-0 name))
  (defstrand card 6 (opts opts) (nr nr) (sel sel-0) (nc nc) (r r) (c c))
  (defstrand reader 5 (sel sel) (nc nc) (opts opts-0) (nr nr) (r r) (c c-0))
  (precedes ((0 3) (1 3)) ((1 2) (0 2)) ((1 4) (0 4)))
  (non-orig (privk r) (privk c))
  (uniq-orig sel-0 nc opts-0 nr)
  (operation encryption-test (added-strand reader 5)
    (enc nr nc (privk r)) (0 4))
  (traces
    ((recv (cat r opts)) (send (cat r c sel-0)) (recv (cat r c nr))
      (send (cat r c nr nc)) (recv (enc nr nc (privk r)))
      (send (enc nr nc (privk c)))))
    ((send (cat r opts-0)) (recv (cat r c-0 sel)) (send (cat r c-0 nr))
      (recv (cat r c-0 nr nc)) (send (enc nr nc (privk r)))))
  (label 1)
  (parent 0)
  (realized)
  (shape)
  (maps ((0) ((r r) (c c) (opts opts) (nr nr) (sel sel-0) (nc nc))))
  (origs (nr (1 2)) (opts-0 (1 0)) (sel-0 (0 1)) (nc (0 3)))))
```

# Translation Goal



**GOAL #1:**  
Highlight the most relevant information.

**Security Goal:**  
Authentication from the Card's pov.

**Must Exist Roles:**  
Reader

**Authenticated Messages:**  
Card Mesg 5 == Reader Mesg 5

**Potential Issues**  
Mesg 1-4 are not authenticated.

**GOAL #2:**  
Provide a suggestion if security goal is not fully satisfied.

**Suggestion:**  
Add the value 'c' to  
Card Mesg 5 and Reader Mesg 5.

Don't need a LLM



Parsing Script

# Translation Goal

ANALYZE IN CPSA

```
(defprotocol blanchet basic
  (defrole init
    (vars (a b name) (s skey) (d data))
    (trace
      (send (enc (enc s b (privk a)) (pubk b)))
      (recv (enc d s))))

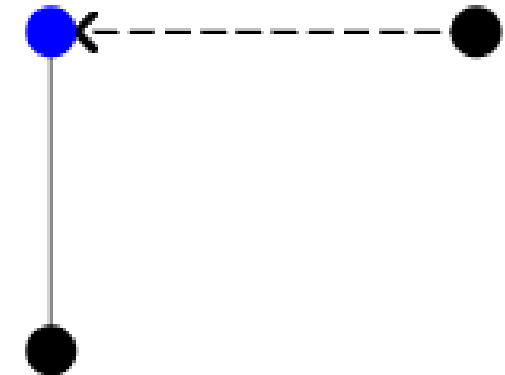
    (defrole resp
      (vars (a b name) (s skey) (d data))
      (trace
        (recv (enc (enc s b (privk a)) (pubk b)))
        (send (enc d s))))

    (comment "Blanchet's protocol using named asymmetric keys"))

  (defskelton blanchet
    (vars (a b name) (s skey) (d data))
    (defstrand resp 2 (a a) (b b) (s s) (d d))
    (non-orig (privk a) (privk b))
    (uniq-orig s)
    (comment "Analyze from the responder's perspective"))
```

resp

init




# Main Conclusions

**Improving the usability for CPSA.  
Combining formal methods with LLMs.**

- **NL → CPSA:** LLMs can reliably serve as translators between NL and CPSA, achieving 98.3% accuracy.
- **CPSA → NL:** LLMs can act as assistants, working in tandem with a CPSA analysis parser to both identify issues in protocols and strengthen them.
- **Future:** More complex protocols, goal translation, CPSA → NL results, integration


# Actual vs Generated: 2\_5 2.scm

```
(defprotocol fpttp_fixed basic
  (defrole init
    (vars (a b p name) (vpa na nb t1 t3 t4 t5 t6
      t7 t8 text) (kab skey)
      (m msg))
    (trace (send (cat vpa b t1 p))
      (recv (cat t4 (enc "9798-2-5_ccf1" vpa kab
        b t3 (ltk a p)) m))
      (send (cat t6 m (enc "9798-2-5_ccf3" na b
```



## Actual CPSA

```
(defprotocol fpttp_fixed basic
  (defrole init
    (vars (a b p name) (na b text) (kab skey)
      (m msg))
    (trace (send (cat "vpa" b "t1" p))
      (recv
        (cat "t4" (enc "9798-2-5_ccf1" "vpa"
          kab b "t3" (ltk a p)) m))
      (send (cat "t6" m (enc "9798-2-5_ccf3" na
```



## Generated CPSA

In CPSA:  
vpa != "vpa"

**Modify Context**

# Actual vs Generated: doorsep 5.scm

## Issue: Incorrect Encryption Key Notations

Allowed CPSA formats:  
(a, invk a) or (pubk a, privk b)  
BUT

Context examples only show  
(pubk a, privk b) form

**Modify Context Examples**

```
(defprotocol doorsep-fixed basic
  (defrole init
    (vars (a b akey) (s skey) (d text))
    (trace
      (send (enc (enc s b (invk a)) b))
      (recv (enc d s))
      (send d))
    (uniq-orig s)))
```

## Actual CPSA

```
(defprotocol doorsep-fixed basic
  (defrole init
    (vars (a b name) (s skey) (d text))
    (trace
      (send (enc (enc s b (invk (pubk a))) (pubk b)))
      (recv (enc d s))
      (send d))
    (uniq-orig s)))
```

## Generated CPSA

# Actual vs Generated: 2-4 2.scm

## Issue: Roles Switched

Logically Equivalent in CPSA  
(CPSA doesn't care about the order of the roles)

**BUT**

Not in the Scoring Method  
(Scoring Method cares about the ordering)

**Modify Scoring Method**

```
protocol threePassMutual_
  (defrole rec
    (vars (a b name)
      (trace
        (recv (cat rb t1))
        (send (cat t3 (e
          a b))))
        (recv (cat t5 (e
          b a))))))
    (uniq-orig ra
```

```
(defrole init
  (vars (a b name)
    (trace
      (send (cat rb t1))
```

**Actual CPSA**

```
protocol threePassMutual_
  (defrole init
    (vars (a b name) (
      (trace
        (send (cat rb t1
          (recv (cat t3 (e
            ltk a b))))
            (send (cat t5 (e
              ltk a b))))))
            (uniq-orig rb))
```

```
(defrole rec
  (vars (a b name) (
    (trace
      (recv (cat rb t1
```

**Generated CPSA**